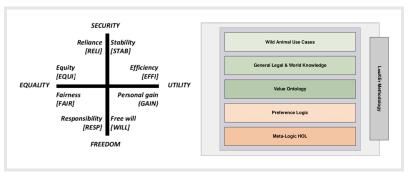
#### Encoding Legal Balancing: Automating an Abstract Ethico-Legal Value Ontology in Preference Logic

#### Christoph Benzmüller, David Fuenmayor, Bertram Lomfeld

Dep. of Mathematics and Computer Science & Dep. of Law Freie Universität Berlin



Workshop "Deontic logic 2020-2030" (adapted from MLR@KR2020 talk)

#### **REASONABLE MACHINES: A RESEARCH MANIFESTO**

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#### ABSTRACT

Future intelligent autonomous systems (IAS) are inevitably deciding on moral and legal questions. e.g. in self-driving cars, health care or humanmachine collaboration. As decision processes in most modern sub-symbolic IAS are hidden. the simple political plea for transparency, accountability and governance falls short. A sound ecosystem of trust requires ways for IAS to autonomously justify their actions, that is, to learn giving and taking reasons for their decisions. Building on social reasoning models in moral psychology and legal philosophy such an idea of \*BEASONABLE MACHINES\* requires novel, hybrid reasoning tools, ethico-legal ontologies and associated argumentation technology. Enabling machines to normative communication creates trust and opens new dimensions of AI application and human-machine interaction.

#### **CORE OBJECTIVES**

- enabling argument-based explanations & justifications of IAS decisions,
- enabling ethico-legal reasoning about, and public critique of, IAS decisions,
- facilitating political and legal governance of IAS decision making,
- evolving ethico-legal agency and communicative capacity of IASs,
- enabling trustworthy human-interaction by normative communication,
- fostering development of novel neuro-symbolic Al architectures.

Long-term vision: To enable machines to give and take normative reasons for their decisions and actions capacitates them to engage in communicative action within social systems.

#### [ARTIFICIAL] SOCIAL REASONING MECHANISM ([A]SRM)

The parallel to human SRM guides the overall architectural design of REASONABLE MACHINES.



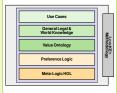
#### REFERENCES

- BENZMÜLLER, C., FUENMAYOR, D., AND LOMFELD, B. Encoding legal balancing: Automating an abstract ethico-legal value ontology in preference logic. In WS on Models of Legal Reasoning, hosted by KR (2020).
- [2] BENZMÜLLER, C., PARENT, X., AND VAN DER TORRE, L. Designing normative theories for ethical and legal reasoning: Logikey framework, methodology, and tool support. Artificial Intelligence 287 (2020), 103348.
- [3] LOMFELD, B. Emotio luris. Skizzen zu einer psychologisch aufgeklärten Methodenlehre des Rechts. In Recht Fühlen (2017), Köhler, Müller-Mall, Schmidt, and Schnädelbach, Eds., Fink, München, pp. 19–32.
- [4] LOMFELD, B. Grammatik der Rechtfertigung: Eine kritische Rekonstruktion der Rechts(fort)bildung. Kritische Justiz 52, 4 (2019).



#### LOGIKEY METHODOLOGY

LOGIKEY [2] is a framework & methodology for the design and engineering of ethical reasoners, normative theories and deontic logics. For recent formalization work see [1] and logikey.org.



#### VALUE BASED EXPLANATIONS

The reflexive symbolic/sub-symbolic feedback loop uses value categories based on an established discoursive moral grammar scheme [4], which we already encoded for legal balancing [1].



#### MODULAR STRUCTURE OF REASONABLE MACHINES RESEARCH

## **Motivation and Contribution**

### **Bigger Vision:**

Reasonable Machines: A Research Manifesto (Benzmüller & Lomfeld, Kl'2020, http://dx.doi.org/10.1007/978-3-030-58285-2\_20)

## Enabling machines to legal balancing

- Challenges: which logic? which value ontology? how to encode? interaction with other legal/world knowledge? which expressivity?
- LogiKEy-Solution: holistic, pluralistic framework; simultaneous modeling at different abstraction layers ... until reflective equilibrium is reached

## Main Contributions:

#### A: Universal (Meta-)Logical Reasoning and LogiKEy approach

- first-time application to support legal balancing
- first-time encoding of preference logic by vanBenthem et al.

### B: Lomfeld's Value Ontology

- first-time operationalization on the computer
- in combination with preference logic by vanBenthem et al.

### C: Combining A&B to model legal balancing in "Wild Animal Cases"

## (A) Universal (Meta-)Logical Reasoning in HOL

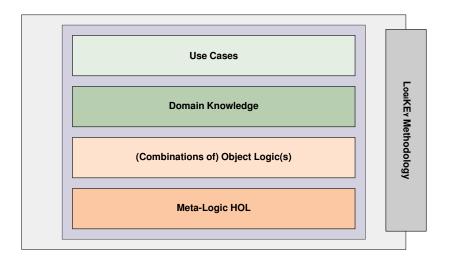
[Science of Computer Programming (2019) vol. 172]



How to Tame the Logic Zoo?

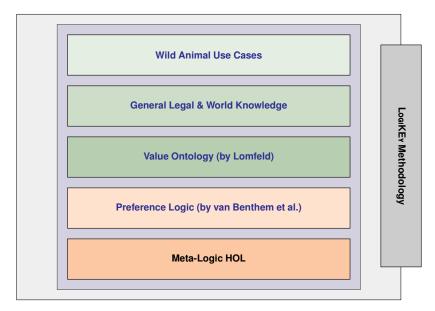
## (A) LogiKEy Methodology

[Artificial Intelligence (2020) vol. 287]



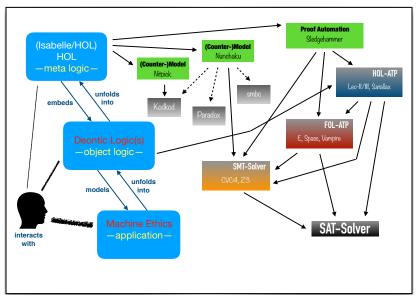
## (A) LogiKEy Methodology

[Artificial Intelligence (2020) vol. 287]



#### (A) Universal (Meta-)Logical Reasoning in Isabelle/HOL

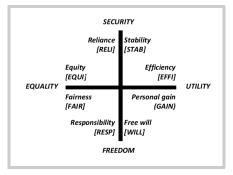
[Science of Computer Programming (2019) vol. 172]



## (B) Value Ontology and Preference Logic

#### Choice of Value Ontology:

Discoursive Grammar of Justification [Lomfeld (2015/2019)]



#### **Choice of Formalization Logic:**

(Modal) Logic for Preferences [vanBenthemGirardRoy(2009), JPL]

#### Der Springer Link

#### Published: 13 August 2008

Everything Else Being Equal: A Modal Logic for *Ceteris Paribus* Preferences

Johan van Benthem 🖂, Patrick Girard & Olivier Roy

Journal of Philosophical Logic 38, 83–125(2009) | Cite this article 393 Accesses | 75 Citations | 0 Altmetric | Metrics

## (B) "Discourse Logic" of Legal Balancing

- Legal reasoning is seen as practical argumentation with a two-level model of (more abstract) values & principles and (more concrete) legal rules.
- Legal rules (or common-law precedents) can be reconstructed as conditional preference relations between conflicting underlying value principles (cf. Alexy 2000; Lomfeld 2015)

**Example:** "In view of events  $E_1$  (a virus pandemic occurs) and  $E_2$  (voluntary shut-down fails) countrywide lock-down becomes sanctioned, since health security outweighs freedom to move."

Application of a rule **R** involves balancing value principles **A** (SECURITY) and **B** (FREEDOM) *in context* (conditions  $E_1$  and  $E_2$ ):

$$R: (E_1 \wedge E_2) \to A > B$$

Acts as justification for the rule's legal consequence (e.g. sanctioned lock-down).

## (B) Encoding using a Logic of Preferences:

#### Choice of Formalization Logic:

[vanBenthemGirardRoy(2009)JPL]

#### D Springer Link

#### Published: 13 August 2008

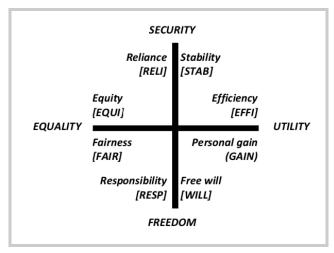
Everything Else Being Equal: A Modal Logic for *Ceteris Paribus* Preferences

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## (B) Value Ontology

But which value principles are to be balanced? [Lomfeld (2015), (2019)]



In our case studies: a decision promoting a particular value (over others) corresponds to ruling for a certain party. (Values are indirectly 'assigned' to particular parties/actors using 'factors'.)

## (B) Value Ontology

# Comparison between some relevant value-based approaches in the literature [Lomfeld (2020)]

VALUES & legal principles	Berman & Hafner 1993	Bench-Capon et al 2005	Bench-Capon 2012	Gordon & Walton 2012	Prakken 2002	Sartor 2002 (Sartor 2010)
<i>FREEDOM</i> - Free choice (WILL) - Responsibility (RESP)	"Protect from interference"	"Court should not make law"	"Reward"			("Liberty")
SECURITY - Stability (STAB) - Reliance (RELI)	"Certainty"	"Clear law"	"Legal certainty" "Public order"	"Security"	"Legal certainty"	("Security") "Less litigation" "Sec. possession"
<i>EQUALITY</i> - Fairness (FAIR) - Equity (EQUI)	"Property rights" "Public land"	"Property"	"Fairness"	"Fairness" "Equity"	"Property rights"	
UTILITY - Efficiency (EFFI) - Personal gain (GAIN)	"Free enterprise and competition"	"Useful" & "Economic activity"	"Utility" - "Econ. valuable" - "Personal gain"		"Economic benefit for society"	"Productivity"

Ambition:

To consistently cover existing value sets from formal argumentation and AI & Law accounts on value-based reasoning, e.g. (Berman and Hafner 1993; Bench-Capon 2012; Gordon and Walton 2012; Sartor 2010).

## (C) Case Study: Pierson vs. Post



Maybe the most famous property law case in American legal history:

Post, a fox hunter, was chasing a fox through public land when Pierson came across the fox and, knowing it was being chased, killed the fox and took it away. Post sued Pierson for damages against his possession of the fox. Post argued that giving chase to the fox was sufficient to establish possession.

## (C) Case Study: Pierson vs Post



- A local court first ruled in favour of Post.
- However, Pierson appealed the ruling to the New York Supreme Court of Judicature, who reversed the decision
- The court ruled in favor of Pierson; citing ancient and modern precedents: "pursuit alone vests no property" (Justinian); and "corporal possession creates legal certainty" (Pufendorf).

## (C) Case Study: Pierson vs Post



In our framework:

- The decision in favour of Pierson implies: STAB(ility) > WILL.
- For "wild animal cases": the legal certainty created by corporal possession (STAB) has preference over "pursuit alone" (WILL).
- Notice the context of validity for the value preference above. Alternatively, Post might argue against this being a "wild animal case".

## (C) Case Study: Conti vs ASPCA



Another famous property law case concerning (wild?) animals:

Chester, a parrot owned by the ASPCA (animal shelter), escaped and was recaptured by Conti. The ASPCA found this out and reclaimed Chester from Conti.

## (C) Case Study: Conti vs ASPCA



In this case, the court ruled in favour of the ASPCA:

- For domestic animals the value preference relation as in Pierson's case does not apply,
- For a domestic animal it is sufficient that the owner did not neglect or stopped caring for the animal, i.e. give up the responsability for its maintenance (RESP).
- This, together with ASPCA's reliance (RELI) in the parrot's property, outweighs Conti's corporal possession (STAB) of the animal.

## Isabelle/HOL Encodings&Tests

## Preference Logic

- Preference Logic Tests
- Value Ontology
- Value Ontology Tests
- General (World) Knowledge
- Pierson Case
- Conti Case

```
1 theory PreferenceLogicBasics (** Benzmüller & Fuenmayor, 2020 **)
     imports Main
 3 begin (** SSE of prefer. logic by van Benthem et al., IPL 2009 **)
 4 (*unimportant*)declare[[syntax ambiguity warning=false]]
    nitpick params[user axioms,expect=genuine,show all,format=3]
 6 (*preliminaries*)
 7 typedecl i
                                    (*possible worlds*)
 8 type synonym \sigma="i\Rightarrowbool" (*'world-lifted' propositions*)
 9 type_synonym \gamma = "i\Rightarrowi\Rightarrowbool" (*preference relations*)
10 type synonym \mu = \sigma \Rightarrow \sigma
                                               (*unary logical connectives*)
11 type synonym \nu = "\sigma \Rightarrow \sigma \Rightarrow \sigma" (*binary logical connectives*)
12 type synonym \pi = "\sigma \Rightarrow bool" (*sets of world-lifted propositions*)
13 (*betterness relation \prec and strict betterness relation \prec*)
14 consts BR::\gamma (" \prec ")
15 abbreviation SBR:: \gamma (" \prec ") where "\vee \prec w \equiv (\vee \prec w) \land \neg (w \prec v)"
16 abbreviation "reflexive R \equiv \forall x, R \times x"
17 abbreviation "transitive R \equiv \forall x \lor z, R x \lor \land R \lor z \longrightarrow R x z"
18 abbreviation "is total R \equiv \forall x y, R x y \lor R y x"
19 axiomatization where rBR: "reflexive BR" and tBR: "transitive BR"
    (*modal logic connectives (operating on truth-sets)*)
abbreviation c1::\sigma (" \perp") where " \perp = \lambda w. False"
22 abbreviation c2::\sigma ("T") where "T = \lambda w. True"
23 abbreviation c3::\mu ("¬ ") where "¬\phi \equiv \lambda w. \neg (\phi w)"
24 abbreviation c4::\nu (infix!" \wedge"85) where "\phi \wedge \psi \equiv \lambda w.(\phi w) \wedge (\psi w)"
    abbreviation c5::\nu (infix!"\vee"83) where "\sigma \vee \psi \equiv \lambda w. (\sigma w) \vee (\psi w)"
26 abbreviation c6::\nu (infix!" \rightarrow "84) where "\phi \rightarrow \psi \equiv \lambda w.(\phi w) \rightarrow (\psi w)"
27 abbreviation c7::\nu (infix!"\leftrightarrow"84) where "\phi \leftrightarrow \psi \equiv \lambda w.(\phi w) \leftrightarrow (\psi w)"
28 abbreviation c8::\mu ("\Box \exists ") where "\Box \exists \varphi \equiv \lambda w. \forall v. (w \prec v) \longrightarrow (\varphi v)"
29 abbreviation c9::\mu ("\diamond \exists ") where "\diamond \exists \varphi \equiv \lambda w. \exists v. (w \prec v) \land (\varphi v)"
  abbreviation c10::\mu ("\Box^{\prec}") where "\Box^{\prec}\varphi \equiv \lambda w.\forall v.(w \prec v) \longrightarrow (\varphi v)"
31 abbreviation cl1::\mu ("\diamond \prec") where "\diamond \neg \varphi \equiv \lambda w. \exists v. (w \prec v) \land (\varphi v)"
32 abbreviation c12::\mu ("E ") where "E\varphi \equiv \lambda w. \exists v. (\varphi v)"
33 abbreviation c13::\mu ("A ") where "A\phi \equiv \lambda w. \forall v. (\phi v)"
34 (*meta-logical predicate for global and validity*)
35 abbreviation g1::\pi ("|_|") where "\psi \equiv \forall w. \psi w"
36 (*some tests: dualities*)
37
    lemma "|(◊≍φ)↔(¬□≍¬φ)| ∧ |(◊≺φ)↔(¬□≺¬φ)| ∧
           |(\mathbf{A}_{\varphi})\leftrightarrow(\neg \mathbf{E}_{\neg \varphi})|^{*} by blast (*proof*)
    (**** Section 3: A basic modal preference language ****)
    (*Definition 5*)
41 abbreviation p1::ν ("_≤εε_")
               where "(\phi \prec_{EE} \psi) u = \exists s. \exists t. \phi s \land \psi t \land s \prec t"
43 abbreviation p2::v (" ≤AE ")
               where "(\varphi \prec_{AE} \psi) u \equiv \forall s. \exists t. \varphi s \longrightarrow (\psi t \land s \prec t)"
```

### Isabelle/HOL Encodings&Tests

### Preference Logic

- Preference Logic Tests
- Value Ontology
- Value Ontology Tests
- General (World) Knowledge
- Pierson Case
- Conti Case

```
1 theory PreferenceLogicBasics (** Benzmüller & Fuenmayor, 2020 **)
         imports Main
     3 begin (** SSE of prefer. logic by van Benthem et al., IPL 2009 **)
     4 (*unimportant*)declare[[syntax ambiguity warning=false]]
     5 nitpick params[user axioms,expect=genuine,show all,format=3]
     6 (*preliminaries*)
     7 typedecl i
                                         (*possible worlds*)
39 (**** Section 3: A basic modal preference language ****)
     (*Definition 5*)
41 abbreviation p1::\nu (" \leq_{EE} ")
                where "(\phi \prec_{FF} \psi) u \equiv \exists s. \exists t. \phi s \land \psi t \land s \prec t"
43 abbreviation p2::v (" ≺AE ")
                where "(\varphi \prec_{AE} \psi) \downarrow \equiv \forall s. \exists t. \varphi s \longrightarrow (\psi t \land s \prec t)"
45 abbreviation p3::ν ("_≺εε_")
                where "(\varphi \prec_{EE} \psi) u = \exists s. \exists t. \varphi s \land \psi t \land s \prec t"
    abbreviation p4::ν (" ≺AE ")
                where "(\varphi \prec_{AE} \psi) u \equiv \forall s. \exists t. \varphi s \longrightarrow (\psi t \land s \prec t)"
49 abbreviation p5:: ν (" ≺AA ")
                where "(\phi \prec_{AA} \psi) u \equiv \forall s. \forall t, (\phi s \land \psi t) \longrightarrow s \prec t"
51 abbreviation p6::ν (" ≻EA ")
52
                where "(\varphi \succ_{EA} \psi) u \equiv \exists s. \forall t. (\varphi s \land \psi t) \longrightarrow t \prec s"
53 abbreviation p7::ν ("_<AA_")
                where "(\varphi \prec_{AA} \psi) u \equiv \forall s. \forall t. (\varphi s \land \psi t) \longrightarrow s \prec t"
55 abbreviation p8::ν (" ≻EA ")
                where "(\varphi \succ_{\mathsf{EA}} \psi) u \equiv \exists s. \forall t. (\varphi s \land \psi t) \longrightarrow t \prec s"
57 abbreviation P1::\nu (" \preceq_{EE} ") where "\varphi \preceq_{EE} \psi \equiv E(\varphi \land \Diamond \exists \psi)"
58 abbreviation P2::\nu (" \prec_{AE} ") where "\phi \prec_{AE} \psi = A(\phi \rightarrow \Diamond \exists \psi)"
59 abbreviation P3::\nu (" \prec_{EE} ") where "\varphi \prec_{EE} \psi \equiv E(\varphi \land \Diamond \neg \psi)"
 60 abbreviation P4::\nu (" \prec_{AE} ") where "\varphi \prec_{AE} \psi \equiv A(\varphi \rightarrow \Diamond \neg \psi)"
61 abbreviation P5::\nu (" \prec_{AA} ") where "\phi \prec_{AA} \psi \equiv \mathbf{A}(\psi \rightarrow \Box \exists \neg \phi)"
62 abbreviation P6::\nu (" \succ_{EA} ") where "\varphi \succ_{EA} \psi \equiv E(\varphi \land \Box \exists \neg \psi)"
63 abbreviation P7::\nu ("\preceq_{AA}") where "\varphi \preceq_{AA} \psi \equiv A(\psi \rightarrow \Box \lnot \neg \varphi)"
64 abbreviation P8::\nu ("_\succeq EA") where "\varphi \succeq EA \psi \equiv E(\varphi \land \Box \lor \neg \psi)"
 65 (*quantification for objects of arbitrary type*)
66 abbreviation mforall ("\forall") where "\forall \Phi \equiv \lambda w, \forall x, (\Phi \times w)"
67 abbreviation mforallB (binder"\forall"[8]9) where "\forall x. \varphi(x) \equiv \forall \varphi"
68 abbreviation mexists ("\exists") where "\exists \Phi \equiv \lambda w. \exists x. (\Phi \times w)"
69 abbreviation mexistsB (binder"3"[8]9) where "\exists x, \varphi(x) \equiv \exists \varphi"
70 (*polymorph operators for sets of worlds/values*)
71 abbreviation subs (infix "\Box" 70) where "A\BoxB = \forall x, A x \longrightarrow B x"
72 abbreviation union (infixr "\sqcup" 70) where "A\sqcupB \equiv \lambda x. A x \lor B x"
73 abbreviation inters (infixr "\square" 70) where "A\squareB \equiv \lambda x. A x \land B x"
74 (*Consistency confirmed (trivial: only abbreviations introduced)*)
75 lemma True nitpick[satisfy user axioms] oops
76 end
```

## Isabelle/HOL Encodings&Tests

- Preference Logic
- Preference Logic Tests
- Value Ontology
- Value Ontology Tests
- General (World) Knowledge
- Pierson Case
- Conti Case

```
theory PreferenceLogicTests1
                                                                (*** Benzmüller & Fuenmavor, 2020 ***)
          imports PreferenceLogicBasics
 3 begin (*Tests for the SSE of van Benthem et al. JPL 2009, in HOL*)
  4 (*Fact 1: definability of the principal operators and verification*)
   lemma F1 9: "(\phi \prec_{FF} \psi) \mathbf{u} \leftrightarrow (\phi \prec_{FF} \psi) \mathbf{u}" by smt
  6 lemma F1 10: "(\phi \prec_{AE} \psi) \downarrow \longleftrightarrow (\phi \prec_{AE} \psi) \downarrow" by smt
  7 lemma F1 11: "(\varphi \prec_{FF} \psi) u \leftrightarrow (\varphi \prec_{FF} \psi) u" by smt
 8 lemma F1 12: "(\varphi \prec_{AE} \psi) \downarrow \longleftrightarrow (\varphi \prec_{AE} \psi) \downarrow" by smt
 9 (*Fact 2: definability of remaining pref. operators and verification*)
10 lemma F2 13: "is total SBR \longrightarrow ((\varphi \prec_{AA} \psi) \cup \longleftrightarrow (\varphi \prec_{AA} \psi) \cup)" by smt
11 lemma F2 14: "is_total SBR \longrightarrow ((\varphi \succ_{EA} \psi) u \longleftrightarrow (\varphi \succ_{EA} \psi) u)" by smt
12 lemma F2 15: "is total SBR \longrightarrow ((\varphi \prec_{AA} \psi) \downarrow \longleftrightarrow (\varphi \prec_{AA} \psi) \downarrow)" by smt
13 lemma F2 16: "is total SBR \longrightarrow ((\varphi \succ_{FA} \psi) \downarrow \longleftrightarrow (\varphi \succ_{FA} \psi) \downarrow)" by smt
14 (*Section 3.5 "Axiomatization" -- verify interaction axioms*)
15 lemma Incl 1: "|(\diamond \neg \varphi) \rightarrow (\diamond \neg \varphi)|"
                                                                             by auto
16 lemma Inter 1: "|(\diamond \exists \diamond \neg \varphi) \rightarrow (\diamond \neg \varphi)|"
                                                                            using tBR by blast
17 lemma Trans le: "|(\diamond \neg \diamond \neg \varphi) \rightarrow (\diamond \neg \varphi)|" using tBR by blast
18 lemma Inter 2: |(\varphi \land \Diamond \exists \psi) \rightarrow ((\Diamond \exists \psi) \lor \Diamond \exists (\psi \land \Diamond \exists \varphi))| by blast
19 Lemma F4: "|(\varphi \land \Diamond \exists \psi) \rightarrow ((\Diamond \exists \psi) \lor \Diamond \exists (\psi \land \Diamond \exists \varphi))| \leftrightarrow
                         (\forall w, \forall v, (((w \prec v) \land \neg (v \prec w)) \longrightarrow (w \prec v)))" by smt
21 lemma Inter_3: "|(\diamond \neg \diamond \neg \varphi) \rightarrow (\diamond \neg \varphi)|" using tBR by blast
22 lemma Incl_2: "|(\diamond \leq \varphi) \rightarrow (E\varphi)|"
                                                                        by blast
23 (*Section 3.6 "A binary preference fragment"*)
24 (* ≺FE is the dual of ≺AA *)
25 lemma "[(\varphi \preceq_{EE} \psi) \leftrightarrow \neg(\psi \prec_{AA} \varphi)] \land [(\varphi \prec_{AA} \psi) \leftrightarrow \neg(\psi \preceq_{EE} \varphi)]" by simp
26 (* ≤<sub>EE</sub> is the dual of ≺<sub>AA</sub> only if totality is assumed*)
27 lemma "(\phi \prec_{EE} \psi) \leftrightarrow \neg(\psi \prec_{AA} \phi)" nitpick oops (*countermodel*)
28 lemma "(\varphi \prec_{EE} \psi) \rightarrow \neg(\psi \prec_{AA} \varphi)" by blast (*this direction holds*)
29 lemma "is total SBR \longrightarrow |(\varphi \prec_{EE} \psi) \leftrightarrow \neg(\psi \prec_{AA} \varphi)|" by blast
30 lemma "(\varphi \prec_{AA} \psi) \leftrightarrow \neg(\psi \preceq_{EE} \varphi) |" nitpick oops (*countermodel*)
31 lemma "(\phi \prec_{AA} \psi) \rightarrow \neg(\psi \prec_{EE} \phi)]" by blast (*this direction holds*)
32 lemma "is_total SBR \longrightarrow |(\varphi \prec_{AA} \psi) \leftrightarrow \neg(\psi \preceq_{EE} \varphi)|" by blast
33 (* verify p.97-98 *)
34 lemma monotonicity: "|((\varphi \leq_{EE} \psi) \land A(\varphi \rightarrow \xi)) \rightarrow (\xi \leq_{EE} \psi)|" by blast
35 lemma reducibility:
                   "|(((\varphi \prec_{\mathsf{EE}} \psi) \land \alpha) \prec_{\mathsf{EE}} \beta) \leftrightarrow ((\varphi \prec_{\mathsf{EE}} \psi) \land (\alpha \prec_{\mathsf{EE}} \beta))|" by blast
37 lemma reflexivity: ||_{\varphi} \rightarrow (\varphi \prec_{\text{EE}} \varphi)|^{*} using rBR by blast
38 (*The condition below is supposed to enforce totality of the preference
          relation. However there are countermodels. See p.98?*)
40 lemma "is total SBR →
                 |((\phi \prec_{\mathsf{EE}} \phi) \land (\psi \prec_{\mathsf{EE}} \psi)) \rightarrow ((\phi \prec_{\mathsf{EE}} \psi) \lor (\psi \prec_{\mathsf{EE}} \phi))|^{"} by auto
42
    Lemma "|((\varphi \prec_{\mathsf{EE}} \varphi) \land (\psi \prec_{\mathsf{EE}} \psi)) \rightarrow ((\varphi \prec_{\mathsf{EE}} \psi) \lor (\psi \prec_{\mathsf{EE}} \varphi))|
                 → is total SBR" nitpick oops (*countermodel - error in paper?*)
44 lemma "is total SBR →
                 \lfloor ((\varphi \preceq_{EE} \varphi) \land (\psi \preceq_{EE} \psi)) \rightarrow ((\varphi \preceq_{EE} \psi) \lor (\psi \preceq_{EE} \varphi)) \rfloor^{*} by auto
    lemma "((\varphi \preceq_{EE} \varphi) \land (\psi \preceq_{EE} \psi)) \rightarrow ((\varphi \preceq_{EE} \psi) \lor (\psi \preceq_{EE} \varphi))
                 is total SBR" nitpick oops (*countermodel - error in paper?*)
48 end
```

## Isabelle/HOL Encodings&Tests

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#### Following [vanBenthemGirardRoy(2009)JPL]

Benzmüller, Fuenmayor, Lomfeld

```
1 theory PreferenceLogicTests2
                                                                                                                                                                   (*** Benzmüller & Fuenmavor, 2020 ***)
                           imports PreferenceLogicCeterisParibus
    3 begin (*** Tests for the SSE of van Benthem et al, JPL 2009 ***)
    4 (**** Section 5: Equality-based Ceteris Paribus Preference Logic ****)
          (*Some tests: dualities*)
      6 lemma "\lfloor (\langle \Gamma \rangle \stackrel{\scriptscriptstyle{\leq}}{\to} \varphi) \leftrightarrow \neg ([\Gamma] \stackrel{\scriptscriptstyle{\leq}}{\to} \neg \varphi) \rfloor" by auto
      7 lemma "|((\Gamma) \leq \varphi) \leftrightarrow \neg([\Gamma] \leq \neg \varphi)|" by auto
    8 lemma "|(\langle \Gamma \rangle \varphi) \leftrightarrow \neg ([\Gamma] \neg \varphi)|"
                                                                                                                                                                by auto
    9 (*Lemma 2*)
   10 lemma lemma2_1: "(\diamond \exists \varphi) w \leftrightarrow (\langle \emptyset \rangle \exists \varphi) w" by auto
 11 lemma lemma2 2: "(\diamond \neg \phi) w \leftrightarrow (\langle \emptyset \rangle \neg \phi) w" by auto
   12 Lemma Lemma2 3: "((E_{\varphi}) \otimes (\langle \emptyset \rangle_{\varphi}) \otimes (\langle A_{\varphi}) \otimes (\langle \emptyset \rangle_{\varphi}) \otimes (\langle \emptyset \otimes (\langle \emptyset \rangle_{\varphi}) \otimes (\langle \emptyset \otimes
 13 (**Axiomatization:**)
 14 (*inclusion and interaction axioms *)
 15 lemma Inc1: "|(\langle \Gamma \rangle \neg \varphi) \rightarrow (\langle \Gamma \rangle \neg \varphi)|" by auto
 16 lemma Inc2: "|(\langle \Gamma \rangle \exists \varphi) \rightarrow (\langle \Gamma \rangle \varphi)|" by auto
 17 Lemma Int3: "\lfloor (\langle \Gamma \rangle \exists (\langle \Gamma \rangle \exists \varphi)) \rightarrow (\langle \Gamma \rangle \exists \varphi) \rfloor" by (meson tBR)
            lemma Int4: "|(\langle \Gamma \rangle \land (\langle \Gamma \rangle \land \varphi)) \rightarrow (\langle \Gamma \rangle \land \varphi)|" by (metis tBR)
 19 lemma Int5: "|(\psi \wedge (\langle \Gamma \rangle \exists \varphi)) \rightarrow ((\langle \Gamma \rangle \exists \varphi) \vee (\langle \Gamma \rangle \exists (\varphi \wedge (\langle \Gamma \rangle \exists \varphi))))|"
                   by (metis rBR)
 21 (*ceteris paribus reflexivity*)
 22 Lemma CetPar6: "\varphi \in \Gamma \longrightarrow |\langle (\Gamma) \varphi \rangle \rightarrow \varphi|"
                                                                                                                                                                                                               by blast
 23 lemma CetPar7: "\varphi \in \Gamma \longrightarrow |((\Gamma) \neg \varphi) \rightarrow \neg \varphi|" by blast
 24 (*monotonicity*)
 25 Lemma CetPar8: "\Gamma \subseteq \Gamma' \longrightarrow |(\langle \Gamma' \rangle \varphi) \rightarrow (\langle \Gamma \rangle \varphi)|" by auto
 26 lemma CetPar9: "\Gamma \subseteq \Gamma' \longrightarrow |(\langle \Gamma' \rangle \exists \varphi) \rightarrow (\langle \Gamma \rangle \exists \varphi)|" by auto
 27 lemma CetPar10: "\Gamma \subseteq \Gamma' \longrightarrow |(\langle \Gamma' \rangle \neg \varphi) \rightarrow (\langle \Gamma \rangle \neg \varphi)|" by auto
 28 (*increase (decrease) of ceteris paribus sets*)
 29 Lemma CetParlia: "(\phi \land (\langle \Gamma \rangle (\alpha \land \phi))) \rightarrow (\langle \Gamma \cup \{\phi\} \rangle \alpha)|"
                                                                                                                                                                                                                                                                                                  by auto
 30 Lemma CetParlib: "|((\neg \varphi) \land (\langle \Gamma \rangle (\alpha \land \neg \varphi))) \rightarrow (\langle \Gamma \cup \{\varphi\} \rangle \alpha)|"
                                                                                                                                                                                                                                                                                                by auto
31 Lemma CetPar12a: "|(\varphi \land (\langle \Gamma \rangle \exists (\alpha \land \varphi))) \rightarrow (\langle \Gamma \cup \{\varphi\} \rangle \exists \alpha)|"
                                                                                                                                                                                                                                                                                                  by auto
 32 lemma CetPar12b: |((\neg \varphi) \land ((\Gamma) \exists (\alpha \land \neg \varphi))) \rightarrow ((\Gamma \cup \{\varphi\}) \exists \alpha)|^{*} by auto
 33 lemma CetPar13a: "(\phi \land ((\Gamma) \land (\alpha \land \phi))) \rightarrow ((\Gamma \cup \{\phi\}) \land \alpha)"
                                                                                                                                                                                                                                                                                                   by auto
 34 Lemma CetPar13b: "|((\neg \varphi) \land (\langle \Gamma \rangle \land (\alpha \land \neg \varphi))) \rightarrow (\langle \Gamma \cup \{\varphi\} \rangle \land \alpha)|" by auto
 35 (*Example 1, Lemma 4, Corollary 1 and Lemma5*)
 36 lemma Ex1: "|(([\Gamma] \exists \varphi) \land (\langle \Gamma \rangle \exists \alpha)) \rightarrow (\langle \Gamma \cup \{\varphi\} \rangle \exists \alpha)|" using rBR by auto
37 Lemma Lemma4: "(\langle \Gamma \rangle \exists \varphi) w \longrightarrow (\exists v. (w \leq v) \land (\varphi v))" by simp
 38 Lemma Corl: "(\langle \Gamma \rangle \varphi) = (\exists v. (w \equiv_{\Gamma} v) \land (\varphi v))" by simp
 39 lemma Lemma5: "(w \triangleleft_{\Gamma} v) \leftrightarrow ((w \prec v) \land (w \equiv_{\Gamma} v))"
                                                                                                                                                                                                                                                            by auto
40 (**** Section 6: Ceteris Paribus Counterparts ****)
 41 (*AA-variant (drawing upon von Wright's)*)
 42 lemma "(\varphi \prec_{AA} \Gamma \psi) u \leftrightarrow (\varphi \prec_{AA} \Gamma \psi) u" nitpick oops (*Ctm*)
 43 lemma "(\phi \prec_{AA}\Gamma \psi) \downarrow \longrightarrow (\phi \prec_{AA}\Gamma \psi) \downarrow" nitpick oops (*Ctm*)
 44 lemma "(\phi \prec_{AA}^{\Gamma} \psi) u \longrightarrow (\phi \prec_{AA}^{\Gamma} \psi) u" by auto
 45 lemma "is total SBR \longrightarrow (\varphi \prec_{AA}^{T} \psi) u \longleftrightarrow (\varphi \prec_{AA}^{T} \psi) u" by smt
 46 lemma "(\phi \preceq_{AA} \psi) u \leftrightarrow (\phi \preceq_{AA} \psi) u" nitpick oops (*Ctm*)
 47 lemma "(\phi \preceq_{AA} \psi) u \longrightarrow (\phi \preceq_{AA} \psi) u" nitpick cops (*Ctm*)
 48 lemma "(\phi \prec_{aa} \Gamma \psi) u \longrightarrow (\phi \prec_{aa} \Gamma \psi) u" by auto
 49 lemma "is total SBR \longrightarrow (\phi \prec_{AA} \Gamma \psi) \downarrow \longleftrightarrow (\phi \prec_{AA} \Gamma \psi) \downarrow" by smt
 50 (*AE-variant*)
51 lemma leAE cp pref: "(\varphi \preceq_{AE} \psi) u \leftrightarrow (\varphi \prec_{AE} \psi) u" by auto
52 lemma leqAE cp pref: "(\varphi \prec_{AE} \psi) u \leftrightarrow (\varphi \prec_{AE} \psi) u" by auto
```

18

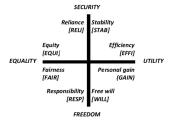
## Isabelle/HOL Encodings&Tests

- Preference Logic
- Preference Logic Tests
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- Conti Case

```
theory PreferenceLogicCeterisParibus (** Benzmüller & Fuenmayor, 2020 **)
        imports PreferenceLogicBasics
    begin (** Ceteris Paribus reasoning by van Benthem et al, JPL 2009 **)
   (*Section 5: Equality-based Ceteris Paribus Preference Logic*)
  6 abbreviation al:: "\sigma \Rightarrow \pi \Rightarrow bool" (" \in ") where "\varphi \in \Gamma \equiv \Gamma \varphi"
  7 abbreviation a2 (" \subseteq ") where "\Gamma \subseteq \Gamma' \equiv \forall \varphi, \varphi \in \Gamma \longrightarrow \varphi \in \Gamma'"
 8 abbreviation a3 (" U ") where "\Gamma U \Gamma' \equiv \lambda \varphi. \varphi \in \Gamma \lor \varphi \in \Gamma'"
 9 abbreviation a4 (" \cap ") where "\Gamma \cap \Gamma' \equiv \lambda \varphi. \varphi \in \Gamma \land \varphi \in \Gamma'"
10 abbreviation a5 ("{ }") where "{\varphi} = \lambda x :: \sigma. x = \varphi"
11 abbreviation a6 ("{ , }") where "{\alpha,\beta} = \lambda x::\sigma. x=\alpha \lor x=\beta"
12 abbreviation a7 ("{,,}") where "{\alpha,\beta,\gamma} \equiv \lambda x::\sigma. x=\alpha \lor x=\beta \lor x=\gamma"
13 abbreviation a8 ("0") where "0 = (\lambda \psi; ;\sigma, False)"
14 abbreviation a9 ("\mathcal{U}") where "\mathcal{U} \equiv (\lambda \psi; ; \sigma, True)"
16 abbreviation cl4 ("_=_") where "w = v \equiv \forall \varphi, \varphi \in \Gamma \longrightarrow (\varphi w \leftrightarrow \varphi v)"
17 abbreviation c15 ("_⊴_") where "w ⊴r v = w ≤ v ∧ w ≡r v"
18 abbreviation c16 ("_<_") where "w <pre>v<<pre>v<<pre>v w v w w w v v
19 abbreviation c17 ("(_)=_") where "(\Gamma)=\varphi \equiv \lambda w. \exists v. w \leq v \land \varphi v"
20 abbreviation c18 ("[_]=_") where "[\Gamma]=\varphi \equiv \lambda w. \forall v. w \leq v \longrightarrow \varphi v"
21 abbreviation c19 ("(_)-=") where "(\Gamma)-\varphi \equiv \lambda w. \exists v. w \triangleleft r v \land \varphi v"
22 abbreviation c20 ("[] \exists ") where "[\Gamma] \exists \varphi \equiv \lambda w, \forall v, w \triangleleft_{\Gamma} v \longrightarrow \varphi v"
23 abbreviation c21 ("()") where "(\Gamma)\varphi \equiv \lambda w. \exists v. w \equiv_{\Gamma} v \land \varphi v"
24 abbreviation c22 ("[]") where "[\Gamma]\varphi \equiv \lambda w. \forall v. w \equiv_{\Gamma} v \longrightarrow \varphi v"
26 (*Section 6: Ceteris Paribus Counterparts of Binary Pref. Statements*)
27 (*operators below not defined in paper; existence is tacitly suggested.
28 AA-variant draws upon von Wright's. AE-variant draws upon Halpern's.*)
29 abbreviation C23 (" <AA- ")
      where "(\varphi \prec_{AA} \psi) u \equiv \forall s. \forall t. \varphi s \land \psi t \longrightarrow s \triangleleft_{\Gamma} t"
31 abbreviation c24 ("____AA-_")
32 where "(\phi \preceq_{AA} \psi) u \equiv \forall s. \forall t. \phi s \land \psi t \longrightarrow s \triangleleft_{\Gamma} t"
33 abbreviation c25 ("_≺AE-_")
34 where "(\phi \prec_{AE} \psi) u \equiv \forall s, \exists t, \phi s \longrightarrow \psi t \land s \triangleleft r t"
35 abbreviation c26 (" ≺AE- ")
36 where "(\phi \prec_{AE} \psi) u \equiv \forall s, \exists t, \phi s \longrightarrow \psi t \land s \triangleleft_{\Gamma} t"
37 abbreviation c27 (" \prec_{AA^{-}}") where "\phi \prec_{AA^{\Gamma}} \psi \equiv A(\psi \rightarrow [\Gamma] \exists \neg \phi)"
38 abbreviation c28 (" \prec_{AA^{-}} ") where "\varphi \prec_{AA^{\Gamma}} \psi \equiv A(\psi \rightarrow [\Gamma] \prec \neg \varphi)"
39 abbreviation c29 (" \prec_{AE^-} ") where "\varphi \prec_{AE^{\Gamma}} \psi \equiv A(\varphi \rightarrow \langle \Gamma \rangle \exists \psi)"
40 abbreviation c30 (" \preceq_{AE^-} ") where "\varphi \preceq_{AE^{\Gamma}} \psi \equiv A(\varphi \rightarrow \langle \Gamma \rangle \exists \psi)"
42 (*Consistency confirmed (trivial: only abbreviations are introduced*)
43 lemma True nitpick[satisfy,user axioms] oops
44 end
```

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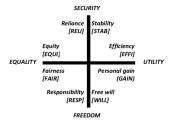
#### Following [Lomfeld(2019)KritischeJustiz]

Benzmüller, Fuenmayor, Lomfeld

```
1 theory ValueOntology
                                   (*** Benzmüller, Fuenmavor & Lomfeld, 2020 ***)
      imports PreferenceLogicBasics
 3 begin (*** Lomfeld's value ontology is encoded ***)
 5 (*two legal parties (there can be more in principle)*)
 6 datatype c = p | d (*parties/contenders: plaintiff, defendant*)
 7 fun other:: "c \Rightarrow c" (" -1") where "p^{-1} = d" | "d^{-1} = p"
 9 consts For:: "c \Rightarrow \sigma^* (*decision: find/rule for party*)
10 axiomatization where ForAx: "[For x \leftrightarrow (\neg For x^{-1})]"
12 datatype (*ethico-legal upper values (wrt, a given party)*)
13 't VAL = FREEDOM 't | UTILITY 't | SECURITY 't | EQUALITY 't
14 type synonym v = "(c)VAL⇒bool" (*principles: sets of upper values*)
15 type synonym cv = "c\Rightarrowv" (*principles are specified wrt. a given party*)
17 abbreviation vset1 ("{]}") where "{\varphi} \equiv \lambda x::(c)VAL. x=\varphi"
18 abbreviation vset2 ("{_,}") where "{\alpha, \beta} = \lambda x::(c)VAL. x=\alpha \lor x=\beta"
20 abbreviation utility::cv ("UTILITY-") where "UTILITY = {UTILITY x}"
21 abbreviation security::cv ("SECURITY-") where "SECURITY = {SECURITY x}"
22 abbreviation equality::cv ("EQUALITY-") where "EQUALITY = {EQUALITY x}"
23 abbreviation freedom::cv ("FREEDOM-") where "FREEDOM* = {FREEDOM *}
24 abbreviation stab::cv ("STAB-") where "STAB* = {SECURITY x, UTILITY x}"
25 abbreviation effi::cv ("EFFI-") where "EFFI* = {UTILITY ×. SECURITY x}"
26 abbreviation gain::cv ("GAIN-") where "GAIN* = {UTILITY ×. FREEDOM ×}"
z7 abbreviation will::cv ("WILL-") where "WILL* = {FREEDOM x. UTILITY x}"
28 abbreviation resp::cv ("RESP-") where "RESP* = {FREEDOM x, EQUALITY x}"
29 abbreviation fair::cv ("FAIR-") where "FAIR* = {EQUALITY x, FREEDOM x}"
   abbreviation equi::cv ("EQUI-") where "EQUI* = {EQUALITY x, SECURITY x}"
31 abbreviation reli::cv ("RELI-") where "RELI× = {SECURITY x, EQUALITY x}"
33 (*derivation operators (cf. theory of "formal concept analysis") *)
34 consts Vrel::"i⇒(c)VAL⇒bool" ("I") (*incidence relation worlds-values*)
35 abbreviation intension::"\sigma \Rightarrow v" ("_^") where "W^ = \lambda v. \forall x. W x \longrightarrow I \times v"
36 abbreviation extension:: "v\Rightarrow \sigma" ("___") where "V__ \equiv \lambda w. \forall x. V x \longrightarrow \mathcal{I} w x"
38 (*shorthand notation for aggregating values*)
abbreviation agg (infixr "\oplus"80) where "v_1 \oplus v_2 \equiv v_1 \square v_2"
40 abbreviation add1 ("[ ]") where "[v] = v1"
41 abbreviation agg2 ("[\oplus]") where "[v_1 \oplus v_2] \equiv (v_1 \oplus v_2)]"
42 abbreviation agg3 ("[_\oplus_\oplus]") where "[v_1 \oplus v_2 \oplus v_3] \equiv (v_1 \oplus v_2 \oplus v_3)]"
43 abbreviation agg4 ("[\oplus \oplus \oplus]") where "[v_1 \oplus v_2 \oplus v_3 \oplus v_4] = (v_1 \oplus v_2 \oplus v_3 \oplus v_4)]"
45 (*chosen variant for preference relation (cf. van Benthem et al. 2009*)
46 abbreviation relPref:: "\sigma \Rightarrow \sigma \Rightarrow \sigma" (" \prec ") where "\varphi \prec \psi \equiv \psi \succ_{EA} \varphi"
47 abbreviation relPrefval:: "v \Rightarrow v \Rightarrow \sigma" (" \prec_v ") where "\varphi \prec_v \psi = \psi \downarrow \succ_{EA} \varphi \downarrow"
49 abbreviation incost ("INCONS_") where (*inconsistency for value support*)
    "INCONS" \equiv [SECURITY"] \land [EQUALITY"] \land [FREEDOM"] \land [UTILITY"]"
                                                                                            18
51
```

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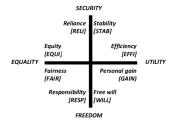
#### Following [Lomfeld(2019)KritischeJustiz]

```
1 theory ValueOntologyTest (*** Benzmüller, Fuenmayor & Lomfeld, 2020 ***)
    imports ValueOntology
 3 begin (* value ontology tests *)
 4 (*values in two opposed quadrants: inconsistent*)
 5 lemma "|[RESP*] ∧ [STAB*] → INCONS*|" by simp
 6 lemma "[RELI*] ∧ [WILL*] → INCONS*|" by simp
 7 (*all values in two non-opposed quadrants: consistent*)
 8 lemma "|[WILL*] ∧ [STAB*] → INCONS*|" nitpick oops (*countermodel*)
 9 (*values in opposed guadrants for different parties: consistent*)
10 lemma "[EQUI*] ∧ [GAIN*] → (INCONS* ∨ INCONS*)|" nitpick oops (*ctm*)
11 Lemma "[RESP*] ∧ [STAB*] → (INCONS* ∨ INCONS*) " nitpick oops (*ctm*)
12 Lemma "[[RELIP] A [WILLP]]" nitpick[satisfy] nitpick oops (*contingent*)
13 (*value preferences tests*)
14 lemma "|WILL× → STAB×| → |WILL× → RELI×⊕STAB×|" by blast
15 lemma "|RELI×⊕STAB× ≺v WILL×| → |STAB× ≺v WILL×|" by auto
16 lemma "|WILL× ≺v RELI×⊕STAB×| → |WILL× ≺v STAB×|"
17 nitpick nitpick[satisfy] oops (*contingent*)
18 lemma "|STAB<sup>×</sup> →<sub>V</sub> WILL<sup>×</sup>| → |RELI<sup>×</sup>⊕STAB<sup>×</sup> →<sub>V</sub> WILL<sup>×</sup>|"
19 nitpick nitpick[satisfy] oops (*contingent*)
28 end
```

```
Nitpick found a model for card i = 1:
  Types:
    c = \{d, p\}
    c VAL =
      {FREEDOM d. FREEDOM p. UTILITY d. UTILITY p.
         EQUALITY d, EQUALITY p, SECURITY d, SECURITY p}
  Constants:
    BR = (\lambda x, ...)((i_1, i_1) := True)
    For = (\lambda x.)((d, i_1) := False, (p, i_1) := True)
    \mathcal{I} = (\lambda \mathbf{x}, \mathbf{x})
         ((i1, FREEDOM d) := False,
          (i1. FREEDOM p) := True.
          (i1, UTILITY d) := False,
          (i1. UTILITY p) := True.
          (i1. EQUALITY d) := False,
          (i1, EOUALITY p) := True,
          (i1. SECURITY d) := False.
          (i1, SECURITY p) := True)
    other = (\lambda x.)(d := p, p := d)
```

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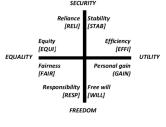
#### Following [Lomfeld(2019)KritischeJustiz]

Benzmüller, Fuenmayor, Lomfeld

```
1 theory ValueOntologyTestLong (** Benzmüller, Fuenmayor & Lomfeld, 2020 **)
 2
      imports ValueOntology
 3 begin
 4 lemma "True" nitpick[satisfy, show all, card i=10] oops
 5 lemma "|INCONSP|" nitpick[satisfy.card i=4] nitpick oops (*contingent*)
 6 (*ext/int operators satisfy main properties of Galois connections*)
 7 lemma G:
                     "B \sqsubseteq A<sup>\uparrow</sup> \longleftrightarrow A \sqsubseteq B<sup>\downarrow</sup>" by blast
 8 lemma G1:
                     "A ⊑ A↑↓" by simp
 9 lemma G2:
                     "B ⊑ B↓↑" by simp
10 lemma G3:
                     "A<sub>1</sub> \sqsubset A<sub>2</sub> \longrightarrow A<sub>2</sub><sup>↑</sup> \sqsubset A<sub>1</sub><sup>↑</sup>" by simp
11 Lemma G4:
                     "B_1 \sqsubseteq B_2 \longrightarrow B_2 \downarrow \sqsubseteq B_1 \downarrow" by simp
12 lemma cll:
                     "A↑ = A↑↓↑" by blast
13 lemma cl2:
                     "B| = B|1| by blast
14 lemma dualla: (A_1 \sqcup A_2)\uparrow = (A_1\uparrow \sqcap A_2\uparrow)^* by blast
   lemma dual1b: (B_1 \sqcup B_2) \downarrow = (B_1 \downarrow \sqcap B_2 \downarrow) by blast
16 lemma
                     (A_1 \sqcap A_2) \uparrow \sqsubseteq (A_1 \uparrow \sqcup A_2 \uparrow) nitpick oops
17 lemma
                     (B_1 \sqcap B_2) \perp \Box (B_1 \sqcup B_2 \perp) nitpick oops
18 lemma dual2a: "(A_1\uparrow \sqcup A_2\uparrow) \sqsubseteq (A_1 \sqcap A_2)\uparrow" by blast
19 lemma dual2b: "(B1↓ ⊔ B2↓) ⊑ (B1 ⊓ B2)↓" by blast
20 (*Note: two different but logically equivalent notations*)
21 lemma "[WILL*] = WILL*1" by simp
22 lemma "[WILL<sup>×</sup>⊕STAB<sup>×</sup>] = (WILL<sup>×</sup>⊕STAB<sup>×</sup>)↓" by simp
24 lemma "|[RELIP] ∧ [WILLP] → INCONSP|" by simp
25 Lemma "INCONSP → [RELIP] ∧ [WILLP]]" by simp
26 Lemma "|[RELIP] A [WILLP]|" nitpick[satisfy] nitpick oops (*contingent*)
27 Lemma "|[FAIR<sup>d</sup>] ∧ [EFFI<sup>d</sup>]|" nitpick[satisfy] nitpick oops (*contingent*)
28 lemma "|(¬INCONS<sup>p</sup>) ∧ [FAIR<sup>d</sup>] ∧ [EFFI<sup>d</sup>]|<sup>*</sup>
29 nitpick[satisfy, show all] nitpick oops (*contingent: p & d independent*)
30 lemma "|(¬INCONS<sup>d</sup>) ∧ (¬INCONS<sup>p</sup>) ∧ [RELI<sup>d</sup>] ∧ [WILL<sup>p</sup>]|"
31 nitpick[satisfy,show all] nitpick oops (*contingent: p & d independent*)
32 (*** more tests ***)
33 (*values in two non-opposed guadrants (nog): consistent*)
  lemma "|[WILL*] ∧ [STAB*] → INCONS*|" nitpick oops (*countermodel found*)
35 Lemma "[WILL*] ∧ [GAIN*] ∧ [EFFI*] ∧ [STAB*] → INCONS*|" nitpick oops
36 (*values in two opposed quadrants: inconsistent*)
37 lemma "|[RESP*] ∧ [STAB*] → INCONS*|" by simp
38 (*values in three quadrants: inconsistent*)
39 Lemma "|[WILL*] ∧ [EFFI*] ∧ [RELI*] → INCONS*|" by simp
40 (*values in opposed quadrants for different parties: consistent*)
41 Lemma "[[EQUI*] ∧ [GAIN*] → (INCONS* ∨ INCONS*)|" nitpick oops (*cntmdl*)
42 Lemma "|[RESP*] ∧ [STAB<sup>y</sup>] → (INCONS<sup>×</sup> ∨ INCONS<sup>y</sup>)|" nitpick cops (*cntmdl*)
43 (*value preferences tests*)
44 lemma "|WILL* → WILL*⊕STAB*|"
     nitpick nitpick[satisfy] oops (*contingent*)
46 lemma "|WILL<sup>×</sup> ≺<sub>V</sub> STAB<sup>×</sup>| → |WILL<sup>×</sup> ≺<sub>V</sub> WILL<sup>×</sup>⊕STAB<sup>×</sup>|" by blast
47 lemma "|WILL<sup>×</sup> ≺<sub>V</sub> STAB<sup>×</sup>| → |WILL<sup>×</sup> ≺<sub>V</sub> RELI<sup>×</sup>⊕STAB<sup>×</sup>|" by blast
48 lemma "|WILL<sup>×</sup> ≺<sub>V</sub> WILL<sup>×</sup>⊕STAB<sup>×</sup>| → |WILL<sup>×</sup> ≺<sub>V</sub> STAB<sup>×</sup>|"
     nitpick nitpick[satisfy] oops (*contingent*)
50 Lemma "|WILL<sup>×</sup> ≺<sub>V</sub> RELI<sup>×</sup>⊕STAB<sup>×</sup>| → |WILL<sup>×</sup> ≺<sub>V</sub> STAB<sup>×</sup>|"
     nitpick nitpick[satisfy] oops (*contingent*)
52 lemma "¬[WILL×⊕STAB× ≺v WILL×]" using rBR by auto
                                                                                                    18
TO TADA RINTLEY OCTADA A MALEY | CTADA A MALEY R MA AND
```

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#### Following [Lomfeld(2019)KritischeJustiz]

Benzmüller, Fuenmayor, Lomfeld

```
1 theory GeneralKnowledge (*Benzmüller, Fuenmayor & Lomfeld, 2020*)
     imports ValueOntology
 3 begin (*** General Legal and World Knowledge (LWK) ***)
 5 (*LWK: kinds of situations addressed*)
   (*appropriation of objects in general*)
 7 consts appAnimal::"σ"
                                (*appropriation of animals in general*)
 8 consts appWildAnimal:: "\sigma" (*appropriation of wild animals*)
 9 consts appDomAnimal:: "\sigma" (*appropriation of domestic animals*)
   (*LWK: postulates for kinds of situations*)
   axiomatization where
    W1: "|(appWildAnimal ∨ appDomAnimal) ↔ appAnimal|" and
    W2: "|appWildAnimal ↔ ¬appDomAnimal|" and
14
15
    W3: "|appWildAnimal → appAnimal|" and
   W4: "|appDomAnimal → appAnimal|" and
16
   W5: "|appAnimal → appObject|"
   (*...further situations regarding appropriation of objects, etc.*)
18
   (*LWK: (prima facie) value preferences for kinds of situations*)
   axiomatization where
22 R1: "|appAnimal \rightarrow (STAB<sup>p</sup> \prec_{v} STAB<sup>d</sup>)|" and
23 R2: "|appWildAnimal → (WILL<sup>x-1</sup> ≺<sub>v</sub> STAB<sup>x</sup>)|" and
24 R3: "|appDomAnimal → (STAB<sup>x-1</sup> ≺<sub>y</sub> RELI<sup>x</sup>⊕RESP<sup>x</sup>)|"
25 (*...further preferences...*)
27 (*LWK: domain vocabulary*)
                (*declares new type for 'entities'*)
28 typedecl e
29 consts Animal::"e \Rightarrow \sigma"
  consts Domestic:: "e \Rightarrow \sigma"
31 consts Fox:: "e \Rightarrow \sigma"
32 consts Parrot::"e \Rightarrow \sigma"
33 consts Pet:: "e \Rightarrow \sigma"
34 consts FreeRoaming:: "e \Rightarrow \sigma"
36 (*LWK: taxonomic (domain) knowledge*)
   axiomatization where
38 W6: "|∀a. Fox a → Animal a|" and
39 W7: "|∀a. Parrot a → Animal a|" and
40 W8: "|∀a. (Animal a ∧ FreeRoaming a ∧ ¬Pet a) → ¬Domestic a|"
41 (*...others...*)
43 (*LWK: legally-relevant, situational 'factors'*)
44 consts Own::"c \Rightarrow \sigma"
                            (*object is owned by party c*)
45 consts Poss::"c \Rightarrow \sigma"
                            (*party c has actual possession of object*)
46 consts Intent:: "c \Rightarrow \sigma" (*party c has intention to possess object*) 18
```

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## Conti Case



#### Following [Lomfeld(2019)KritischeJustiz]

```
27 (*LWK: domain vocabulary*)
28 typedecl e
                    (*declares new type for 'entities'*)
29 consts Animal:: "e \Rightarrow \sigma"
   consts Domestic:: "e \Rightarrow \sigma"
31 consts Fox:: "e \Rightarrow \sigma"
32 consts Parrot::"e \Rightarrow \sigma"
33 consts Pet::"e \Rightarrow \sigma"
34 consts FreeRoaming::"e⇒σ"
36 (*LWK: taxonomic (domain) knowledge*)
   axiomatization where
38 W6: "∀a. Fox a → Animal a " and
39 W7: "|∀a. Parrot a → Animal a|" and
    W8: "|∀a. (Animal a ∧ FreeRoaming a ∧ ¬Pet a) → ¬Domestic a|"
40
41
   (*...others...*)
42
43 (*LWK: legally-relevant, situational 'factors'*)
44 consts Own::"c \Rightarrow \sigma"
                              (*object is owned by party c*)
45 consts Poss:: "c \Rightarrow \sigma"
                              (*party c has actual possession of object*)
46 consts Intent:: c \Rightarrow \sigma (*party c has intention to possess object*)
47 consts Mal:: "c \Rightarrow \sigma" (*party c acts out of malice*)
48 consts Mtn:: "c \Rightarrow \sigma" (*party c respons. for maintenance of object*)
50 (*LWK: meaning postulates for general notions*)
51 axiomatization where
52 W9: "|Poss x \rightarrow (\neg Poss x^{-1})|" and
53 W10: "|Own x → (¬Own x<sup>-1</sup>)|"
54 (*...others...*)
56 (*LWK: conditional value preferences, e.g. from precedents*)
57 axiomatization where
58 R4: "(Mal x^{-1} \land 0wn x) \rightarrow (STAB^{x-1} \prec_{v} RESP^{*} \oplus RELI^{*})"
59 (*...others...*)
61 (*LWK: relate values, outcomes and situational 'factors'*)
62 axiomatization where
63 F1: "|For x → (Intent x ↔ □≤[WILL×])|" and
64 F2: "|For x \rightarrow (Mal x^{-1} \leftrightarrow \Box^{\perp}[RESP^{x}])|" and
65 F3: "|For x \rightarrow (Poss x \leftrightarrow \Box \leq [STAB^{\times}])|" and
66 F4: "|For x \rightarrow (Mtn \ x \leftrightarrow \Box = [RESP^{\times}])|" and
67 F5: "|For x \rightarrow (0wn \ x \leftrightarrow \Box^{(RELI^{*})})"
68
69 (*theory is consistent, (non-trivial) model found*)
70 lemma True nitpick[satisfy,card i=10] oops
71 end
```

### Isabelle/HOL Encodings&Tests

- Preference Logic
- Preference Logic Tests
- Value Ontology
- Value Ontology Tests
- General (World) Knowledge
- Pierson Case
- Conti Case



#### Following [Lomfeld(2019)KritischeJustiz]

Benzmüller, Fuenmayor, Lomfeld

```
(*** Benzmüller, Fuenmayor & Lomfeld, 2020 ***)
 1 theory Pierson
     imports GeneralKnowledge
 3 begin (*** Pierson v. Post "wild animal" case **)
 5 (*case-specific 'world-vocabulary'*)
 6 consts α:: "e" (*appropriated animal (fox in this case) *)
 7 consts Pursue:: "c \Rightarrow e \Rightarrow \sigma"
 8 consts Capture:: "c \Rightarrow e \Rightarrow \sigma"
10 (*case-specific taxonomic (legal domain) knowledge*)
   axiomatization where
12 CW1: "|(\exists c. Capture c \alpha \land \neg Domestic \alpha) \rightarrow appWildAnimal|" and
   CW2: "|\forall c. Pursue c \alpha \rightarrow Intent c|" and
   CW3: "|\forall c. Capture c \alpha \rightarrow Poss c|"
16 lemma True nitpick[satisfy,card i=4] oops (*satisfiable*)
   18
19 abbreviation "Pierson facts \equiv |Fox \alpha \wedge (FreeRoaming \alpha) \wedge
     (\neg \text{Pet } \alpha) \land \text{Pursue p } \alpha \land (\neg \text{Pursue d } \alpha) \land \text{Capture d } \alpha|^*
21
22 (*decision for defendant (Pierson) is compatible with premises*)
23 Lemma "Pierson facts ∧ |¬INCONS<sup>P</sup>| ∧ |¬INCONS<sup>d</sup>| ∧ |For p ≺ For d|"
     nitpick[satisfy.card i=4] oops (* (non-trivial) model found*)
24
26 (*decision for plaintiff (Post) is compatible with premises*)
27 lemma "Pierson_facts ∧ |¬INCONS<sup>P</sup>| ∧ |¬INCONS<sup>d</sup>| ∧ |For d ≺ For p|"
     nitpick[satisfy.card i=4] oops (* (non-trivial) model found*)
29
30 (*decision for defendant (Pierson) is provable*)
31 theorem assumes Pierson_facts shows "|For p ≺ For d|"
32
     by (metis assms CW1 CW2 W6 W8 ForAx R2 F1 other.simps(2) rBR)
34 (*while a decision for the plaintiff is not*)
35 lemma assumes Pierson facts shows "|For d ≺ For p|"
36 nitpick[card i=4] oops (*counterexample found*)
   39 (* Theory amendment: the animal is not free-roaming since it
      is being chased by a professional hunter (Post) *)
   consts Hunter:: "c \Rightarrow \sigma"
42 axiomatization where (*case-specific legal rule for hunters*)
43 R5: "|(Hunter x \land Pursue x \alpha) \rightarrow (STAB<sup>x-1</sup> \prec_v EFFI<sup>x</sup>)|"
45 abbreviation "Post facts \equiv |Fox \alpha \land (¬FreeRoaming \alpha) \land
                                                                             18
```

Hunter p  $\land$  Pursue p  $\alpha$   $\land$  (¬Pursue d  $\alpha$ )  $\land$  Capture d  $\alpha$ |"

46

### Isabelle/HOL Encodings&Tests

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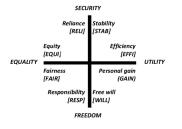


#### Following [Lomfeld(2019)KritischeJustiz]

```
abbreviation rieson_facts = rox \alpha / (rieeRodining \alpha) /
    (\neg \text{Pet } \alpha) \land \text{Pursue p } \alpha \land (\neg \text{Pursue d } \alpha) \land \text{Capture d } \alpha|^*
20
21
   (*decision for defendant (Pierson) is compatible with premises*)
22
   lemma "Pierson facts ∧ |¬INCONSP| ∧ |¬INCONS<sup>d</sup>| ∧ |For p ≺ For d|"
     nitpick[satisfy,card i=4] oops (* (non-trivial) model found*)
24
26
   (*decision for plaintiff (Post) is compatible with premises*)
27
   lemma "Pierson_facts ∧ [¬INCONS<sup>p</sup>] ∧ [¬INCONS<sup>d</sup>] ∧ |For d ≺ For p|"
     nitpick[satisfy,card i=4] oops (* (non-trivial) model found*)
28
29
   (*decision for defendant (Pierson) is provable*)
зө
31 theorem assumes Pierson facts shows "|For p ≺ For d|"
     by (metis assms CW1 CW2 W6 W8 ForAx R2 F1 other.simps(2) rBR)
32
33
34 (*while a decision for the plaintiff is not*)
   lemma assumes Pierson facts shows "|For d ≺ For p|"
36
     nitpick[card i=4] oops (*counterexample found*)
38
   39 (* Theory amendment: the animal is not free-roaming since it
      is being chased by a professional hunter (Post) *)
41 consts Hunter:: "c \Rightarrow \sigma"
   axiomatization where (*case-specific legal rule for hunters*)
43 R5: "(Hunter x \land Pursue x \alpha) \rightarrow (STAB<sup>x-1</sup> \prec_v EFFI<sup>x</sup>)|"
45 abbreviation "Post facts \equiv |Fox \alpha \land (\negFreeRoaming \alpha) \land
      Hunter p \land Pursue p \alpha \land (¬Pursue d \alpha) \land Capture d \alpha|"
46
48 (*decision for defendant (Pierson) is compatible with premises*)
49 lemma "Post facts ∧ |¬INCONS<sup>p</sup>| ∧ |¬INCONS<sup>d</sup>| ∧ |For p ≺ For d|"
     nitpick[satisfy,card i=4] oops (* (non-trivial) model found*)
52
   (*decision for plaintiff (Post) is compatible with premises too*)
53 Lemma "Post facts ∧ |¬INCONSP| ∧ |¬INCONS<sup>d</sup>| ∧ |For d ≺ For p|"
     nitpick[satisfy.card i=4] oops (* (non-trivial) model found*)
54
56 (*indeed, a decision for plaintiff (Post) now becomes provable*)
57 theorem assumes Post facts shows "|For d ≺ For p|"
58 using assms by (metis CW3 ForAx R5 F3 other.simps rBR)
59
60 (*while a decision for the defendant is now refutable*)
61 lemma assumes Post facts shows "|For p ≺ For d|"
    nitpick[card i=4] oops (* counterexample found*)
63 end
```

## Isabelle/HOL Encodings&Tests

- Preference Logic
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#### Following [Lomfeld(2019)KritischeJustiz]

#### (\*\*\* Benzmüller, Fuenmayor & Lomfeld, 2020 \*\*\*) 1 theory Conti imports GeneralKnowledge 3 begin (\*\*\* ASPCA v. Conti "wild animal" case \*\*) (\*case-specific 'world-vocabulary'\*) 6 consts α::"e" (\*appropriated animal (parrot in this case) \*) consts Care:: " $c \Rightarrow e \Rightarrow \sigma$ " 8 consts Prop:: " $c \Rightarrow e \Rightarrow \sigma$ " 9 consts Capture:: " $c \Rightarrow e \Rightarrow \sigma$ " 11 (\*case-specific taxonomic (legal domain) knowledge\*) 12 axiomatization where CW1: "Animal $\alpha \land \text{Pet } \alpha \rightarrow \text{Domestic } \alpha$ " and 13 CW2: "|( $\exists c. Capture \ c \ \alpha \ \land Domestic \ \alpha$ ) $\rightarrow appDomAnimal|$ " and CW3: " $|\forall c. Care c \alpha \rightarrow Mtn c|$ " and 16 CW4: " $\forall c. Prop \ c \ \alpha \rightarrow 0 wn \ c$ ]" and CW5: " $|\forall c.$ Capture $c \alpha \rightarrow Poss c|$ " 18 19 lemma True nitpick[satisfy,card i=4] oops (\*satisfiable\*) abbreviation "ASPCA facts $\equiv$ |Parrot $\alpha \land$ Pet $\alpha \land$ Care p $\alpha \land$ 23 Prop p $\alpha \land (\neg \text{Prop d } \alpha) \land \text{Capture d } \alpha|^*$ 24 (\* decision for defendant (Conti) is compatible with premises\*) 26 Lemma "ASPCA facts ∧ |¬INCONS<sup>p</sup>| ∧ |¬INCONS<sup>d</sup>| ∧ |For p ≺ For d|" 27 nitpick[satisfy.card i=4] oops (\* (non-trivial) model found\*) 28 29 (\* decision for plaintiff (ASPCA) is compatible with premises\*) 30 lemma "ASPCA\_facts ∧ |¬INCONS<sup>p</sup>| ∧ |¬INCONS<sup>d</sup>| ∧ |For d ≺ For p|" nitpick[satisfy,card i=4] oops (\* (non-trivial) model found\*) 32 33 (\* decision for plaintiff (ASPCA) is provable\*) lemma aux: assumes ASPCA facts shows "|(STAB<sup>d</sup> ≺v RELI<sup>p</sup>⊕RESP<sup>p</sup>)|" 34 using CW1 CW2 W7 assms R3 by fastforce 36 theorem assumes ASPCA facts shows "|For d ≺ For p|" 37 using assms aux CW5 ForAx F3 other.simps(1) rBR by metis 39 (\* while a decision for the defendant is refutable\*) 40 lemma assumes ASPCA facts shows "|For p ≺ For d|" nitpick[card i=4] oops (\*(non-trivial) counterexample found\*) 42 end

## Models and Countermodels are particularly helpful!

```
Skolem constant:
 \lambda y, y = (\lambda x, -)(1) := 14, 12 := 11, 13 := 11, 14 := 14)
c = {d, p}
e \times i [boxed] = {(e_1, i_1), (e_1, i_2), (e_1, i_3), (e_2, i_4)}
 c VAL = {FREEDON d, FREEDON p, UTILITY d, UTILITY p, EQUALITY d, EQUALITY p, SECURITY d, SECURITY p}
 Capture =
   (Ax. )
   (d, e1, 11) := True, (d, e1, 12) := True, (d, e1, 13) := True, (d, e1, 14) := True, (p, e1, 11) := False, (p, e1, 12) := False,
    (p, e1, i3) := False, (p, e1, i4) := False)
 Care =
   (\lambda x, \cdot)
   ((d, e1, i1) := False, (d, e1, i2) := False, (d, e1, i3) := False, (d, e1, i4) := True, (p, e1, i1) := True, (p, e1, i2) := True,
    (p, e1, 13) := True, (p, e1, 14) := True)
 Prop =
   ((d, e1, 11) := False, (d, e1, 12) := False, (d, e1, 13) := False, (d, e1, 14) := False, (p, e1, 11) := True, (p, e1, 12) := True,
    (p, e1, in) := True, (p, e1, i4) := True)
 Animal = (\lambda x. )((e1, i1) := True, (e1, i2) := True, (e1, i3) := True, (e1, i4) := True)
 Domestic = (\lambda x. _)({e<sub>1</sub>, i<sub>2</sub>} := True, (e<sub>1</sub>, i<sub>2</sub>) := True, (e<sub>1</sub>, i<sub>3</sub>) := True, (e<sub>1</sub>, i<sub>4</sub>) := True)
 Fox = (λx, )((e1, 11) := False, (e1, 12) := False, (e1, 12) := False, (e1, 14) := False)
 FreeRoaming = (Ax. )({e1, i1} := False, (e1, i2) := False, (e1, i3) := False, (e1, i4) := False)
 Intent =
   (Ax. )
   (d, i1) := False, (d, 12) := True, (d, 13) := False, (d, 14) := True, (p, 11) := False, (p, 12) := False, (p, 13) := True,
    (p. is) := False)
 Liv -
   (Ax. )
   ((d, i1) := False, (d, i2) := True, (d, i3) := False, (d, i4) := False, (p, i1) := False, (p, i2) := False, (p, i3) := False,
    (p, 14) := False)
 Mtn =
   (Ax. )
   ((d, i)) := True, (d, i)) := False, (d, i) := False, (d, i4) := True, (p, i) := True, (p, i) := True, (p, i) := True,
    (p, is) := True)
 Own =
   (\lambda x. __)
   ((d, 1) := False, (d, 1) := False, (d, 1) := False, (d, 1) := False, (p, 1) := True, (p, 1) := True, (p, 1) := True,
    (p, 14) := True)
 Parrot = (\lambda x. ) ((e1, i1) := True, (e1, i2) := True, (e1, i3) := True, (e1, i4) := True)
 Pet = (λx. _)((e1, i1) := True, (e1, i2) := True, (e1, i3) := True, (e1, i4) := True)
 Poss =
   (Ax. )
   ((d, i1) := True, (d, i2) := True, (d, i3) := True, (d, i4) := True, (p, i1) := False, (p, i2) := False, (p, i2) := False,
    (p, is) := False)
  appAnimal = (\lambda x.) (11 := True, 12 := True, 13 := True, 14 := True)
  appDomAnimal = (Ax. )(i1 := True, i2 := True, i3 := True, i4 := True)
 appOblect = (\lambda x, \cdot)(1) := True, 1) := True, 1) := True, 14 := True)
 appWildAnimal = (\lambda x, \cdot)(i_1 := False, i_2 := False, i_3 := False, i_4 := False)
 BR = (\lambda x.)
       ((i1, i1) := True, (i1, i2) := False, (i1, i3) := False, (i1, i4) := False, (i2, i1) := True, (i2, i2) := True,
        (i2, i3) := False, (i2, i4) := False, (i3, i1) := True, (i3, i2) := True, (i3, i3) := True, (i3, i4) := False,
        (14, 11) := False, (14, 12) := False, (14, 13) := False, (14, 14) := True)
 For =
   (\lambda x, 1)
   ((d, i1) := False, (d, i2) := True, (d, i3) := True, (d, i4) := False, (p, i1) := True, (p, i2) := False, (p, i3) := False,
    (p, is) := True)
 I = (\lambda x. )
      ((i1, FREEDOM d) := True, (i1, FREEDOM p) := True, (i1, UTILITY d) := True, (i1, UTILITY p) := False, (i1, EOUALITY d) := False,
       (i1, EQUALITY p) := True, (i1, SECURITY d) := True, (i1, SECURITY p) := True, (i2, FREEDOM d) := True, (i2, FREEDOM p) := True,
       (i2, UTILITY d) := True, (i2, UTILITY p) := False, (i2, EOUALITY d) := True, (i2, EOUALITY p) := False, (i2, SECURITY d) := True,
       (12, SECURITY p) := True, (13, FREEDOM d) := False, (13, FREEDOM p) := True, (13, UTILITY d) := True, (13, UTILITY p) := False,
       (13, EQUALITY d) := True, (13, EQUALITY p) := True, (13, SECURITY d) := True, (13, SECURITY p) := False, (14, FREEDOM d) := True,
       (i4, FREEDOM p) := True, (i4, UTILITY d) := False, (i4, UTILITY p) := False, (i4, EQUALITY d) := True, (i4, EQUALITY p) := True,
       (14. SECURITY d) := True, (14. SECURITY p) := True)
 other = \{\lambda x, \dots\} \{d := p, p := d\}
```

## **Conclusion and Related Work**

### **Contributions:**

- Feasibility study for legal balancing on the computer
- Embedding of Preference Logic in HOL
- Demonstrated formalization&use of Lomfeld's value ontology
- Successful application of
  - LogiKEY methodology and
  - Universal (Meta-)Logical Reasoning in HOL
- Flexibility, Expressiveness and ready to use ATP Support!

#### **Related work:**

- Constructive interpretation in law, including model of value balancing: [Maranhão&Sartor(2019)ICAIL]
- Models to quantify legal balancing: [Alexy(2003), Sartor(2010)]

#### **Bigger Vision:**

#### **Reasonable Machines: A Research Manifesto**

(Benzmüller & Lomfeld, Kl'2020, http://dx.doi.org/10.1007/978-3-030-58285-2\_20)